

# KERNEL.DATA.STRUCTURES

## Task Structure

The fundamental entity of any operating system is the task or a process. The task structure is defined in the **linux/sched.h** and it is the important data structure and following are some of its parameters.

```
struct task_struct {
```

```
volatile long state; //determines the current state of the process like unrunnable,  
runnable and stopped
```

```
unsigned long flags; // various flags for accounting purpose like PF_STARTIN,  
PF_MEMALLOC
```

```
unsigned long ptrace; // process is being monitored by other process.
```

```
int sigpending; //define when the signals must be handed over to this process
```

```
mm_segment_t addr_limit;
```

```
struct exec_domain *exec_domain; //domain other than x86
```

```
long need_resched; //flag indicates that scheduling must be executed
```

```
int lock_depth; //structure is protected before simultaneous access can take place
```

```
long counter; //dynamic priority of a process
```

```
long nice; //static priority of a process
```

```

unsigned long policy; //scheduling policy like SCHED_RR, SCHED_FIFO, SCHED_OTHER
unsigned long rt_priority;
.
};

```

## Process Relations

This data structure describes the relation between processes like child process, parent process, etc.

```
struct task_struct *p_opptr; //original parent process
```

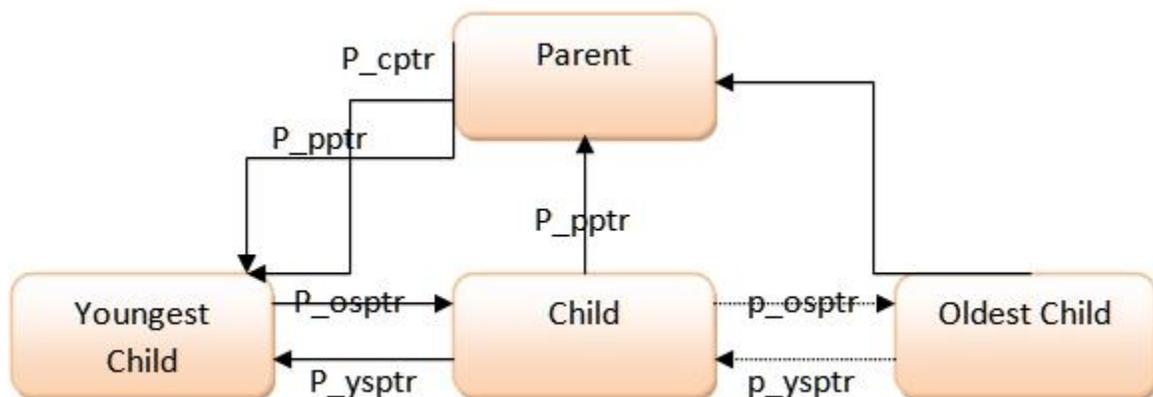
```
struct task_struct *p_pptr; //Parent Process
```

```
struct task_struct *p_cptr; //child process
```

```
struct task_struct *p_ysptr; //youngest sibling process
```

```
struct task_struct *p_osptr; //old sibling process
```

the following diagram depicts the relationship between the process.



## Process ID

Every process has its own process ID called pid defined as follows:

```
pid_t pid, pgrp, session, tgid; //pgrp - process group, session in which the process runs
and tgid is thread group id
```

**int leader;** //there will be a leader process among the process

apart from the above, there are two other parameters like **uid** and **gid** referred namely as user identification and group identification, which enables to create the access right.

## **Timing**

Times are measured in ticks. usually the timing signals are generated by the hardware timer every 10ms which is captured by the timer interrupt.

there are different timing parameters are specified like

**long per\_cpu\_utime[NO OF CPUS];** //user mode

**long per\_cpu\_stime[NO OF CPUS];** //system mode

**struct tms times;** //the timing values are added for all the child processes and other processes

**unsigned long start\_time;** //time at which the process was generated.

## **FILES and INODES**

Files and inodes represent the file structures.

Files defines the attributes of file like **mode of operation, position of the cursor, flags, reference count, directory entry, etc**

inodes defines the information about the files **like owner of the file, device in which the file is located, access rights of the file, size of the file, last access time, last modification time, etc**

**struct file {**

**mode\_t f\_mode;** // mode of operation like read, write, append, etc

**loff\_t f\_pos;** //position of the read/write cursor

**atomic\_t f\_count;** //reference count like how many times the file has been opened by the system call

**unsigned int f\_flags;** //access control

**struct dentry \*fs\_dentry;** // file descriptor

```
.  
};  
  
struct inode {  
  
kdev_t i_dev; // device id in which the file is located  
  
unsigned long i_ino; //file within device  
  
umode_t i_mode; //mode of operation  
  
uid_t i_uid; //user id which is the owner of the file  
  
gid_t i_gid; //group id, the file can be accessed by a group  
  
off_t i_size; //size of the file in bytes  
  
time_t i_mtime; //last modification time  
  
time_t i_atime; //last accessed time  
  
.  
};
```

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