INTERRUPT HARDWARE

We pointed out that an I/O device requests an interrupt by activating a bus line called interrupt-request. Most computers are likely to have several I/O devices that can request an interrupt. A single interrupt-request line may be used to serve n devices as depicted. All devices are connected to the line via switches to ground. To request an interrupt, a device closes its associated switch. Thus, if all interrupt-request signals \( \text{INTR}_1 \) to \( \text{INTR}_n \) are inactive, that is, if all switches are open, the voltage on the interrupt-request line will be equal to \( V_{dd} \). This is the inactive state of the line. Since the closing of one or more switches will cause the line voltage to drop to 0, the value of \( \text{INTR} \) is the logical OR of the requests from individual devices, that is,

\[
\text{INTR} = \text{INTR}_1 + \ldots + \text{INTR}_n
\]

It is customary to use the complemented form, \( \overline{\text{INTR}} \), to name the interrupt-request signal on the common line, because this signal is active when in the low-voltage state.

ENABLING AND DISABLING INTERRUPTS:

The facilities provided in a computer must give the programmer complete control over the events that take place during program execution. The arrival of an interrupt request from an external device causes the processor to suspend the execution of one program and start the execution of another. Because interrupts can arrive at any time, they may alter the sequence of events from the envisaged by the programmer. Hence, the interruption of program execution must be carefully controlled.

Let us consider in detail the specific case of a single interrupt request from one device. When a device activates the interrupt-request signal, it keeps this signal activated until it learns that the processor has accepted its request. This means that the interrupt-request signal will be active during execution of the interrupt-service routine, perhaps until an instruction is reached that accesses the device in question.

The first possibility is to have the processor hardware ignore the interrupt-request line until the execution of the first instruction of the interrupt-service routine has been completed. Then, by using an Interrupt-disable instruction as the first instruction in the
interrupt-service routine, the programmer can ensure that no further interruptions will occur until an Interrupt-enable instruction is executed. Typically, the Interrupt-enable instruction will be the last instruction in the interrupt-service routine before the Return-from-interrupt instruction. The processor must guarantee that execution of the Return-from-interrupt instruction is completed before further interruption can occur.

The second option, which is suitable for a simple processor with only one interrupt-request line, is to have the processor automatically disable interrupts before starting the execution of the interrupt-service routine. After saving the contents of the PC and the processor status register (PS) on the stack, the processor performs the equivalent of executing an Interrupt-disable instruction. It is often the case that one bit in the PS register, called Interrupt-enable, indicates whether interrupts are enabled.

In the third option, the processor has a special interrupt-request line for which the interrupt-handling circuit responds only to the leading edge of the signal. Such a line is said to be edge-triggered.

Before proceeding to study more complex aspects of interrupts, let us summarize the sequence of events involved in handling an interrupt request from a single device. Assuming that interrupts are enabled, the following is a typical scenario.

1. The device raises an interrupt request.
2. The processor interrupts the program currently being executed.
3. Interrupts are disabled by changing the control bits in the PS (except in the case of edge-triggered interrupts).
4. The device is informed that its request has been recognized, and in response, it deactivates the interrupt-request signal.
5. The action requested by the interrupt is performed by the interrupt-service routine.
6. Interrupts are enabled and execution of the interrupted program is resumed.