FIBRE CHANNEL (FC) - 2

FC-2: data transfer

FC-2 is the most comprehensive layer in the Fibre Channel protocol stack. It determines how larger data units (for example, a file) are transmitted via the Fibre Channel network. It regulates the flow control that ensures that the transmitter only sends the data at a speed that the receiver can process it. And it defines various service classes that are tailored to the requirements of various applications.

Exchange, sequence and frame

FC-2 introduces a three-layer hierarchy for the transmission of data (Figure 3.13). At the top layer a so-called exchange defines a logical communication connection between two end devices. For example, each process that reads and writes data could be assigned its own exchange. End devices (servers and storage devices) can simultaneously maintain several exchange relationships, even between the same ports. Different exchanges help the FC-2 layer to deliver the incoming data quickly and efficiently to the correct receiver in the higher protocol layer (FC-3).



Figure 4.13 One sequence is transferred after another within an exchange. Large sequences are broken down into several frames prior to transmission. On the receiver side, a sequence is not delivered to the next highest protocol layer (FC-3) until all the frames of the sequence have arrived.

A sequence is a larger data unit that is transferred from a transmitter to a receiver. Only one sequence can be transferred after another within an exchange. FC-2 guarantees that sequences are delivered to the receiver in the same order they were sent from the transmitter; hence the name 'sequence'. Furthermore, sequences are only delivered to the next protocol layer up when all frames of the sequence have arrived at the receiver (Figure 3.13). A sequence could

represent the writing of a file or an individual database transaction. A Fibre Channel network transmits control frames and data frames. Control frames contain no useful data, they signal events such as the successful delivery of a data frame. Data frames transmit up to 2,112 bytes of useful data. Larger sequences therefore have to be broken down into several frames. Although it is theoretically possible to agree upon different maximum frame sizes, this is hardly ever done in practice.

A Fibre Channel frame consists of a header, useful data (payload) and a Cyclic Redundancy Checksum (CRC) (Figure 3.14). In addition, the frame is bracketed by an start-of-frame (SOF) delimiter and an end-of-frame (EOF) delimiter. Finally, six filling words must be transmitted by means of a link between two frames. In contrast to Ethernet

and TCP/IP, Fibre Channel is an integrated whole: the layers of the Fibre Channel protocol stack are so well harmonised with one another that the ratio of payload to protocol overhead is very efficient at up to 98%. The CRC checking procedure is designed to recognise all transmission errors if the underlying medium does not exceed the specified error rate of 10–12. Error correction takes place at sequence level: if a frame of a sequence is wrongly transmitted, the entire sequence is re-transmitted. At gigabit speed it is more efficient to resend a complete sequence than to extend the Fibre Channel hardware so that individual

lost frames can be resent and inserted in the correct position. The underlying protocol layer must maintain the specified maximum error rate of 10–12 so that this procedure is efficient.



Figure 4.14 The Fibre Channel frame format.

Frame Destination Address (D_ID) Frame Source Address (S ID) Sequence ID Number of the frame within the sequence Exchange ID

Flow control

Flow control ensures that the transmitter only sends data at a speed that the receiver can receive it. Fibre Channel uses the so-called credit model for this. Each credit represents the capacity of the receiver to receive a Fibre Channel frame. If the receiver awards the transmitter a credit of '4', the transmitter may only send the receiver four frames. The transmitter may not send further frames until the receiver has acknowledged the receipt of at least some of the transmitted frames.

FC-2 defines two different mechanisms for flow control: end-to-end flow control and link flow control (Figure 3.15). In end-to-end flow control two end devices negotiate the end-to-end credit before the exchange of data. The end-to-end flow control is realized on the HBA cards of the end devices. By contrast, link flow control takes place at each physical connection. This is achieved by two communicating ports negotiating the buffer-to-buffer credit. This means that the link flow control also takes place at the Fibre Channel switches.

Service classes

The Fibre Channel standard defines six different service classes for exchange of data between end devices. Three of these defined classes (Class 1, Class 2 and Class 3) are realised in products available on the market, with hardly any products providing the connection-oriented Class 1. Almost all new Fibre Channel products (HBAs, switches, storage devices) support the service classes Class 2 and Class 3, which realise a packet-oriented service (datagram service). In addition, Class F serves for the data exchange between the switches within a fabric. Class 1 defines a connection-oriented communication connection between two node ports: a Class 1 connection is opened before the transmission of frames. This specifies a route through the Fibre Channel network. Thereafter, all frames take the same route through the Fibre Channel network so that frames are delivered in the sequence in which they were transmitted. A Class 1 connection guarantees the availability of the full bandwidth. A port thus cannot send any other frames while a Class 1 connection is open.



Figure 4.15 In link flow control the ports negotiate the buffer-to-buffer credit at each link (1). By contrast, in end-to-end flow control the end-to-end credit is only negotiated between the end devices (2).

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