Abstract - This paper presents a method to find the minimum number of phasor measurement units (PMUs) for complete observability of power systems network for normal operating conditions. A linear algorithm is used to determine the minimum number of PMUs needed to make the system observable. For state estimation and fault diagnosis in power system synchronized snapshot of the entire system must be necessary. The proposed method is used to benchmark the optimal PMUs placement solution for the IEEE 14-bus, IEEE 18-bus, IEEE 30-bus and IEEE 57-bus test systems.

Keywords - Phasor Measurement Units (PMUs), Linear Algorithm (LA), Optimal Placement of PMUs (OPP), Global Positioning System (GPS), Power System Observability.

I. INTRODUCTION

Phasor Measurement Units (PMUs) become more and more important and attractive to power engineers because they can provide synchronized measurements of real-time phasors of voltage and currents [1]. As the state estimator plays an important role in the security of power system to enhance state estimation in a problem needed to be solved. Several algorithms have been published in the literature. Until recently, it was not possible to measure phase angles of the bus voltages in real time due to the technical difficulties in synchronizing measurements from distant locations. But introducing the PMUs in power systems, it is possible to measure the real-time phasors of voltages and currents at widely dispersed locations with respect to a global positioning system (GPS) clock [2].

The methodology is needed to determine the optimal location of PMUs in a power system. In addition to its ability to measure voltage and current phasors, a state-of-the-art PMU may include other features such as protective actions. The objective of the present work is to find the minimum number of PMUs to make the system topologically observable, as well as the optimal locations of these PMUs. In recent years, there has been a significant research activity on this problem. In [3], a bisectioning search method is implemented to find the minimum number of PMUs to make the system observable. The simulated annealing method used is used to randomly choose the placement sets to the test for observability at each step of the bisection search. In [1], the authors use a simulated annealing technique in their graph-theoretic procedure to find the optimal PMU locations.

In [5] and [6] the authors use integer programming to determine the minimum number of PMUs. The method, however, may suffer from the problem of being trapped in local minima. Multiple objectives, such as minimizing the measurement redundancy, cannot be handled by integer programming. In [7]-[8], the OPP optimization problem is solved using PSAT, a MATLAB based toolbox, and depth-first search (DeFS) method is compared with other methods. Another depth-first search (DeFS) method is proposed in [9]. The DeFS algorithm is computationally faster, but the solution is not optimum, because the optimization criterion is strict. A modified depth-first approach is the minimum spanning tree (MST) method [9]. The MST algorithm improves the DeFS approach, which also has fast computing characteristics, and improves DeFS’s complex and weak convergence. A novel topological method based on the augment incidence matrix and Tabu Search(TS) algorithm, is proposed in [10]. The solution of the combinatorial OPP problem requires less computation and is highly robust. The method is faster and more convenient than conventional observability analysis methods using complicated matrix analysis, because it manipulates integer numbers. A TS method on meter placement to maximize topological observability is presented in [10]. The GA method suggested in [11] solves the OPP problem using different PMU placement criteria, such as the absence of critical measurements and critical sets from the system, maximum quantity of measurements received as compared to the initial one, maximum accuracy of estimates, minimum cost of PMU placement, and transformation of the network graph into a tree. The immune algorithm (IA) is a search strategy based on genetic algorithm principles and inspired by protection mechanisms of living organisms against bacteria and viruses. In reference [12], the application of the immune genetic algorithm (IGA) method to the OPP problem is presented. Utilization of the local and prior knowledge associated with the considered problem is the main idea behind IGA.
prior knowledge of the OPP problem was inferred based on the topological observability analysis and was abstracted as some vaccines. The injection of these vaccines into the individuals of generations, revealed a remarkable increase in the convergence process. A BPSO algorithm, with the objective of minimum PMU installation costs, is introduced in [13]. A hybrid algorithm based on BPSO and immune mechanism is introduced in [14]. It provides a speedy and general analyzing method of power network topology observation based on the properties of PMU and topological structure information of the power network. The classical ant colony optimization (ACO) algorithm is a probabilistic technique for solving computational problems which can be reduced to finding good paths through graphs. A generalized ACO algorithm is proposed in [15]. The present paper proposes the Linear Algorithm method to minimize the PMUs.

II. PHASOR MEASUREMENT UNIT (PMU) TECHNOLOGY

A phasor measurement unit is a device that provides as a minimum, synchrophasor and frequency measurements for one or more three phase AC voltage and/or current waveforms [16]. These measure are marked with a GPS time stamp in time intervals down to 20 ms [1]. This same time sampling of voltage and current waveforms using a common synchronizing signal from the global positioning set elite ensure synchronicity among PMUs. This synchronicity makes the PMU one of the most important devices for power system control and monitoring.

Fig. 1: PMU Layout with GPS time stamped Signal

Fig. 1 shows PMUs geographically dispersed to form a wide area monitoring system (WAMS) in which the PMUs deliver GPS time-tagged measurements to a Phasor Data Concentrator (PDC). The PDC sorts the incoming phasor measurements before signal processing converts PMU data into actionable information that can be presented to an operator in the form of a Human Machine Interface (HMI). This HMI provides an operator with critical information about the state of power system.

III. MATERIALS AND METHODS

Observability analysis is a fundamental component of real-time state estimation. There are two major algorithms for power network observability analysis: topology-based algorithms and numerical methods. Topology methods are developed from graph theories, compared to numerical methods that are mainly based on numerical factorization of measurement Jacobian matrices. Numerical methods are less suitable for large systems because they are involved with large dimension matrices that increase the computational complexity [1]. If we model buses in a power system by vertices and model the transmission and distribution line connecting buses by edge, this problem is converted to be an optimization problem and requires the extension of the topological observation theory. The observation rules [17] can be described as following:-

Rule 1: A bus with PMU installation is observable, and its adjacent buses are all observable because their voltages can be calculated by Ohm’s law with the help of the PMU measurement.

Rule 2: If a bus is adjacent to an observable zero-injection bus to which all other adjacent buses are observable, then the bus is observable because its voltage can be calculated by KCL and Ohm’s law.

Rule 3: If all buses adjacent to a zero-injection bus is observable because its voltage can be calculated by KCL and Ohm’s law.

To get the fast solution, a good initial guess of PMUs placement, this algorithm was tested for a list of distribution system and proven very good efficient. In 2002, Haynes et al [4], mathematically proved that, for a tree having $k$ vertices of degree at least 3, the “power dominating number”

$$\gamma_p(T) \geq \frac{(k+2)}{3} \quad (1)$$

$$\gamma_p(T) \leq \frac{n}{3} \quad (2)$$

Where $n$ is the total no. of vertices.

Equation (1) and (2) give the upper and lower bounds for the power dominating number. Although a power system does not have to be a tree topology, these theorems corresponded to the computation result from [3] very well. Haynes et al in this paper also gave an algorithm to find out the dominating set $S$ and a partition of the whole set $G$ into $S$ such that each subset induces a “spider”. This algorithm was strictly proven in this paper. In 2009, Mohammadi-Ivatloo summarized most available topological based formulated algorithms, including genetic algorithm (2009), Tabu search (2006), Integer linear programming (2008). However, complexity of each algorithm is still left to be discussed. In engineering practice, we are most interested in the dominating set, i.e., where to mount the PMUs. The partition is not the primary concern. Taking advantage of the upper and lower bounds [4], a linear algorithm is proposed in this study; this algorithm is proven especially effective for small system.
IV. THE LINEAR ALGORITHM

Creating a Node –Incidence Matrix for Power System:

By the topological information of a power system the interconnection of the various buses can be grouped in an array called node-incidence matrix.

Fig.1: IEEE-14 bus system

To produce the node incidence matrix the rule is simple: If node i is connected to node j, then \( A_{ij} = 1 \), where \( i \) and \( j \) and \( A_{ii} = 0 \). Normally \( A \) is a large sparse matrix. For example, for the IEEE 14-bus system,

\[
A = \begin{bmatrix}
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
\end{bmatrix}
\]

It is easy to find out that there are 8 nodes with degree 3 or more. Here \( K=8 \) and \( n=14 \), so, the number of PMUs needed that is , the only possible values for \( S \) between 3 and 4 using equation (1) and (2). Now we want to find out the minimum number of PMUs, and the dominating set \( S \). The basic idea of this algorithm is to test all possible node combinations by the observation rules, until one combination is found to be able to “observe” all the system. We call a test for a combination as a measurement. For the IEEE 14-bus system, the maximum number of measurements is number of combinations produce by selecting numbers of a group in between 3-4, who will converge, will give the required number of PMUs in the system. That is 70 for IEEE 14-bus system. We need to keep in mind that, in the implementation of the algorithm, we may not have to run all the 70 measurements to find out the \( S \)-set. The number of measurement before we get an \( S \)-set (which is usually not unique) can be any number between 1 and 70.

V. ALGORITHM FOR FINDING THE MINIMUM NUMBER OF PMU’S

Begin
1. Read in node-incidence matrix \( A \) with all buses (nodes) in the system says \( G \).
2. Calculate the bounds of \( S \).
3. Check the observability of the system by creating loop starting from the lower bound, to the upper bound:
   a. Generate a node combination, e.g., \{2, 4, 5\}. These nodes are mounted with PMUs, thus observed. Save them in array \( O \).
   b. Find out all nodes adjacent to these 3 or 4 nodes.
   c. Save them in array \( O \).
4. Find out all nodes that are not in \( O \).
   a. Pick up such a node \( j \), use rule 4 to judge if it is observed. If yes, put \( j \) in \( O \) and pick up another node and check.
   b. If all “not-in-\( O \)” nodes have been checked, compare \( O \) to the whole set \( G \).
5. If \( O=G \), output the node combination. That is the \( S \)-set. Quit. If \( O \) is not equal to \( G \), generate another node combination.
6. If solution does not converge then increase the numbers in the group by one toward upper bound.
7. Output how many number of combination has been tried that is the number of measurement.
END

VI. RESULTS AND DISCUSSION

The linear algorithm is implemented by Matlab (version 7.8.0.347) and is tested for IEEE 14, 18, 39, 57-bus systems. The results are given and compared to other methods. Table 1 gives the comparison of computation results by different algorithms. Table 2 gives the bus locations where the PMUs will be installed. Due to large computational time and complexity, it not works properly on large bus network like on 57-bus system and above. The best results are found on IEEE 14, 18, 30-bus systems. The proposed method is very fast in providing optimal solution as compared to other search methods. Simulation results for different network show the effectiveness of the proposed method in finding the minimum optimal number of PMU bus locations for complete observability assessment of Power Systems.

<table>
<thead>
<tr>
<th>Sr. No.</th>
<th>Power System (bus no. N)</th>
<th>Method</th>
<th>V</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>IEEE-14 bus</td>
<td>Proposed Linear Algorithm</td>
<td>3</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Dual Search</td>
<td>3</td>
</tr>
</tbody>
</table>
TABLE I
Comparison of Result with different Algorithm Where N: number of buses , V: minimum number of pmus required

<table>
<thead>
<tr>
<th>Power System</th>
<th>No. of PMU’s</th>
<th>Bus Locations of PMUs</th>
</tr>
</thead>
<tbody>
<tr>
<td>IEEE-14 bus</td>
<td>3</td>
<td>2,6,9</td>
</tr>
<tr>
<td>IEEE-18 bus</td>
<td>4</td>
<td>4,6,12,15</td>
</tr>
<tr>
<td>IEEE-30 bus</td>
<td>7</td>
<td>2,6,10,12,15,25,17</td>
</tr>
<tr>
<td>IEEE-57 bus</td>
<td>N/A</td>
<td>N/A</td>
</tr>
</tbody>
</table>

VII. CONCLUSION

The proposed linear algorithm takes advantage of the upper and lower bounds and the graph theorems that were mathematically proven in [4], which greatly reduced the computation in seeking a dominating set in a power system. Compared to the algorithm in [3] and [18], the proposed algorithm can theoretically guarantee the minimum number of PMUs. Compared to the algorithms in [1], it is simpler and easier to be implemented. However, the computation complexity reflected by measurement study indicates that the linear algorithm should be very competitive to other topology based algorithm and other numerical methods and provides complete observability for the distribution system[1]. For further study we can try to find such solutions which give lesser capital investment for the PMU placement as the number of communication port, environmental concerns, technological issues and life cycle also contribute to the cost/unit of a PMU in the system.

REFERENCES

[17] Chi Su, Zhe Chen Institute of Energy Technology “Optimal Placement of Phasor Measurement Units with New Consideration”.